CONCORDIA UNIVERSITY

DEPARTMENT OF COMPUTER SCIENCE & SOFTWARE ENGINEERING

COMP 335/4	Theoretical Computer Science	Winter 2015
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Assignment 3

Due date: March 27, 2015, by 23:59 p.m. EDT

- 1. In each case, what language is generated by CFG's below. Justify your claim (prove it!)
 - (a) $S \to aSa|bSb|aAb|bAa, A \to aAa|bAb|a|b|\epsilon$
 - (b) $S \to aS|bS|a$
 - (c) $S \rightarrow SS|bS|a$
 - (d) $S \to SaS|b \ S \to aT|bT|\epsilon, \ T \to aS|bS.$
- 2. Find a CFG for each of the languages below.
 - (a) $L = \{a^{n}b^{m} : n \neq m-1\}$ (b) $L = \{a^{n}b^{m}c^{k} : n = m \text{ or } m \neq k\}$ (c) $L = \{w \in \{a, b\}^{*} : n_{a}(w) \neq n_{b}(w)\}$ (d) \overline{L} , where $L = \{w \in \{a, b\}^{*} : w = a^{n}b^{n}, n \geq 0\}$
- 3. In each case below, show that the grammar is ambiguous, and find an equivalent unambiguous grammar.
 - (a) $S \to SS|ab|a$
 - (b) $S \to ABA, A \to aA|\epsilon, B \to bB|\epsilon$
 - (c) $S \to aSb|aaSb|\epsilon$
- 4. Design a PDA to accept each of the following languages. You may design your PDA to accept either by final state or empty stack, whichever is more convenient.
 - (a) The set of strings over $\{0, 1\}$ such that no prefix has more 1's than 0's.
 - (b) The set of strings with twice as many 0's as 1's.
 - (c) The set of strings over $\{a, b\}$ that are *not* of the form ww, that is, not equal to any string repeated.

- 5. Construct a PDA corresponding to the context-free grammar $S \rightarrow SS \mid \{SX \mid [SY \mid \epsilon X \rightarrow \} Y \rightarrow]$ Note that $\{, [,], \text{ and } \}$ are terminals.
- 6. Consider the PDA $P = \{\{q_0, q_1, q_2\}, \{a\}, \{\clubsuit, Z_0\}, \delta, q_0, Z_0, \{q_2\}\}, \text{ where } \delta(q_0, a, Z_0) = \{(q_1, \clubsuit Z_0)\}, \delta(q_1, a, \clubsuit) = \{(q_0, \epsilon)\}, \text{ and } \delta(q_0, \epsilon, Z_0) = \{(q_2, \epsilon)\}.$ Construct a CFG (using the method in the text) corresponding to P.
- 7. Use the Pumping Lemma for CFL's to show that none of the following languages are context-free.
 - (a) $L_1 = \{ww : w \in \{a, b\}^*\}$ (b) $L_2 = \{a^n b^k : 0 \le n \le k^2\}$ (c) $L_3 = \{a^n b^m c^k : 0 \le n < m, n \le k \le m\}$
- 8. Convert the following grammar into Chomsky normal form

$$\begin{split} S &\to aA | aBB \\ A &\to aaA | \epsilon \\ B &\to bB | bbC \\ C &\to C | B \end{split}$$

9. (a) Show that the language

$$L = \{a^n b^n : a, b \in \{a, b\}, n \text{ is not a multiple of 5} \}$$

is context-free.

(b) Let $L = \{a^n b^n : n \ge 0\}$, and $M = \{a^{2m} b^{2p} : m \ge 0, p \ge 0\}$. Construct a PDA for L and a DFA¹ for M. Then use the Cartesian construction to obtain a PDA for $L \cap M$.

¹Leave out the trap state