## Concordia University

Introduction to Theoretical Computer Science Winter 2015

## Solution to Assignment 2 - Part II

5. If $L$ is regular, it is accepted by some DFA, say $A=\left(Q, \Sigma, \delta, s_{0}, F\right)$. We will construct an $\epsilon$-NFA, such that $L(B)=\operatorname{third}(L(A))$. Here you need four copies of $A$. Formally,

$$
B=(Q \times\{1,2,3,4\}, \Sigma, \rho, F \times\{2,3,4\}),
$$

where $\rho=$
$\{(\langle p, 1\rangle, \epsilon,\langle q, 2\rangle):(p, a, q) \in \delta$, for some $a \in \Sigma\} \cup$
$\{(\langle p, 2\rangle, \epsilon,\langle q, 3\rangle):(p, a, q) \in \delta$, for some $a \in \Sigma\} \cup$
$\{(\langle p, 3\rangle, \epsilon,\langle q, 4\rangle):(p, a, q) \in \delta\} \cup$
$\{(\langle p, 4\rangle, \epsilon,\langle p, 1\rangle): p \in Q\}$
6. The following DFA describes the language $L$.


Call this DFA $A$. Clearly $L(A)=L$. More precisely, $A=\left(\left\{r_{0}, r_{1}, r_{2}\right\},\{0,1\}, \delta,\left\{r_{1}, r_{2}\right\}\right)$. Construct $B=\left(\left\{r_{0}, r_{1}, r_{2}\right\},\{0,1\}, \gamma,\left\{r_{1}, r_{2}\right\}\right)$, where

$$
\gamma(q, a)=\hat{\delta}(q, h(a))
$$

for all $q \in\left\{r_{1}, r_{2}, r_{3}\right\}$ and all $a \in\{0,1\}$.
Then $L(B)=h^{-1}(L(A))=h^{-1}(L)$. The transition diagram for $B$ is shown below.

7. Table of distinguishabilities:


- Using the table we can construct the minimum state equivalent DFA.


8. (a) $\left\{a^{n} b^{m}: n \geq 2, m \geq 1\right\}$


To show that this automaton is minimal, we compute its table of distinguishabilities:

| B | $\times$ |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
| C | $\times$ | $\times$ |  |  |
| D | $\times$ | $\times$ | $\times$ |  |
| E | $\times$ | $\times$ | $\times$ | $\times$ |
|  | A | B | C | D |

(b) $\left\{a^{n} b: n \geq 0\right\} \cup\left\{a b^{n}: n \geq 1\right\}$


To show that this automaton is minimal, we compute its table of distinguishabilities:

| B | $\times$ |  |  |  |  |
| :---: | :---: | :---: | :---: | :---: | :---: |
| C | $\times$ | $\times$ |  |  |  |
| D | $\times$ | $\times$ | $\times$ |  |  |
| E | $\times$ | $\times$ | $\times$ | $\times$ |  |
| F | $\times$ | $\times$ | $\times$ | $\times$ | $\times$ |
|  | A | B | C | D | E |

(c) $\left\{a^{n} b: n \geq 0\right\} \cup\left\{a b^{n}: n \geq 1\right\}$


To show that this automaton is minimal, we compute its table of distinguishabilities:

| B | $\times$ |  |  |  |
| :---: | :---: | :---: | :---: | :---: |
| C | $\times$ | $\times$ |  |  |
| D | $\times$ | $\times$ | $\times$ |  |
| E | $\times$ | $\times$ | $\times$ | $\times$ |
|  | A | B | C | D |

